

# MonteBoxFinder: Detecting and Filtering Primitives to Fit a Noisy Point Cloud

## Supplementary Material

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Project page <https://michaelramamonjisoa.github.io/projects/MonteBoxFinder>

## 1 Supplementary Video

In addition to this supplementary document, we provide a supplementary video showing progress through iterations for our method and compared baselines.

## 2 More on Cuboids

### 2.1 Constructing Cuboids from Pairs of Plane Segments

In this section, we provide more details regarding the construction of cuboids given a pair of plane segments  $\pi_A = (X_A, \mathbf{N}_A)$  and  $\pi_B = (X_B, \mathbf{N}_B)$ .

**Checking planes adjacency** We recall that two planes  $(\pi_A, \pi_B)$  should be considered for constructing a cuboid only if they fulfill two requirements: *alignment* and *proximity*.

*Proximity* Two plane segments are adjacent if they verify the *proximity* criterion, which requires that they have at least one connected component, such that  $\min(\text{ChamferDistance}(X_A \rightarrow X_B), \text{ChamferDistance}(X_B \rightarrow X_A)) < \gamma$ , where  $\gamma$  is a small 3D distance.

*Alignment* Two plane segments  $\pi_A = (X_A, \mathbf{N}_A)$  and  $\pi_B = (X_B, \mathbf{N}_B)$  are aligned if they are either “orthogonal enough” or “co-linear enough”; This corresponds to  $|\mathbf{N}_A^T \mathbf{N}_B| < \alpha$  or  $|\mathbf{N}_A^T \mathbf{N}_B| > \beta$ , respectively, where  $\alpha \ll 1$  and  $\beta \lesssim 1$ . In our experiments we set  $(\alpha, \beta, \gamma) = (0.3, 0.7, 0.05m)$ .

**Getting cuboids main axes** We can now construct *two* orthonormal bases  $\mathcal{B}_A = (\mathbf{u}_A, \mathbf{v}_A, \mathbf{w}_A)$  and  $\mathcal{B}_B = (\mathbf{u}_B, \mathbf{v}_B, \mathbf{w}_B)$  of vectors using Gram-Schmidt orthonormalization with  $N_A$  or  $N_B$  as the first vector alternatively, as shown in Equations (1a) and (1b).

$$\begin{aligned}
\mathbf{u}_A &:= \frac{\mathbf{N}_A}{\|\mathbf{N}_A\|_2} \\
\mathbf{v}_A &:= \mathbf{N}_B - \frac{\mathbf{u}_A^T \mathbf{N}_B}{\|\mathbf{N}_B\|_2} \mathbf{u}_A \quad (1a) \\
\mathbf{v}_A &\leftarrow \frac{\mathbf{v}_A}{\|\mathbf{v}_A\|_2} \\
\mathbf{w}_A &:= \mathbf{u}_A \times \mathbf{v}_A,
\end{aligned}$$

$$\begin{aligned}
\mathbf{u}_B &:= \frac{\mathbf{N}_B}{\|\mathbf{N}_B\|_2} \\
\mathbf{v}_B &:= \mathbf{N}_A - \frac{\mathbf{u}_B^T \mathbf{N}_A}{\|\mathbf{N}_A\|_2} \mathbf{u}_B \quad (1b) \\
\mathbf{v}_B &\leftarrow \frac{\mathbf{v}_B}{\|\mathbf{v}_B\|_2} \\
\mathbf{w}_B &:= \mathbf{u}_B \times \mathbf{v}_B,
\end{aligned}$$

**Computing the final cuboids** Based on the two cuboids bases, we compute their sizes by simply projecting all 3D points  $X \in X_A \cup X_B$  and computing the minimum and maximum of the projections along each axes of  $\mathcal{B}_A$  and  $\mathcal{B}_B$ . This results in two bounding boxes aligned with  $\mathcal{B}_A$  and  $\mathcal{B}_B$  respectively, which both enclose all points in  $X_A \cup X_B$ .

## 2.2 Computing Intersections with `isCompatible`

In order to check compatibility between cuboids  $s_1$  and  $s_2$ , we design a variation of an Intersection-over-Union criterion, replacing the *Union* with the volume of the smallest cuboid between  $s_1$  and  $s_2$ . In order to compute the volume of the intersection, we approximate volumes by sampling points in both  $s_1$  and  $s_2$  and counting points that are inside both. Full details of the procedure used in `isCompatible` are given in Algorithm 1. In practice we use an intersection threshold  $\eta = 10\%$ .

## 3 More About our Baselines

### 3.1 The Hill-Climbing Algorithm

The Hill-Climbing algorithm [1] is a naive greedy descent algorithm that constructs a solution iteratively, where at each iteration, it comprehensively searches for the proposal that best improves the objective function of a solution  $\mathcal{S}_F$ , while leaving the solution valid *i.e.* with no incompatibilities. If no proposal is available nor can improve the objective function, the algorithm stops  $\times$ . A pseudo-code for the Hill-Climbing algorithm is provided in Algorithm 2.

*Sub-optimal solutions* In practice, Hill-Climbing leads to sub-optimal solutions where the algorithm gets stuck into a local minimum. This is because the algorithm first greedily fits large regions of the scene, therefore employing large `Cuboid`; This makes a lot of potentially good cuboids unavailable, as they would intersect with that large `Cuboid`.



**Algorithm 1:** The isCompatible function

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procedure isCompatible( $s_1, \mathcal{S}, \eta$ )
  Result: Returns True if Cuboid  $s_1$  is compatible with all Cuboid in  $\mathcal{S}$ 
  Input      : Cuboids Cuboid  $\mathcal{S}$ , threshold  $\eta$ 
  if ( $\forall s_2 \in \mathcal{S}, \text{IntersectionOverVolume}(s_1, s_2) > \eta$ ) then
    | return False ;
  return True ;

procedure IntersectionOverVolume( $s_1, s_2$ )
  Input      : Cuboids  $s_1$  and  $s_2$ 
  Volume of Cuboid  $s_1$   $V_1 := \text{Volume}(s_1)$ 
  Volume of Cuboid  $s_2$   $V_2 := \text{Volume}(s_2)$ 
  Number of samples  $N_{\text{samples}} := 5000$ ;
  Number of samples from  $s_1$  in  $s_2$   $N_{1 \subset 2} := 0$ ;
  Number of samples from  $s_2$  in  $s_1$   $N_{2 \subset 1} := 0$ ;
  // Sample 3D points within both cuboids  $s_1$  and  $s_2$ 
   $\mathcal{X}_1 := \text{sample\_points.inside}(s_1, N)$  ;
   $\mathcal{X}_2 := \text{sample\_points.inside}(s_2, N)$  ;
  // Count points sampled in  $s_1$  which are also inside  $s_2$ 
  for ( $x \in \mathcal{X}_1$ ) {
    | if  $x \in s_2$  then
    | |  $N_{1 \subset 2} \leftarrow N_{1 \subset 2} + 1$ ;
  }
  // Count points sampled in  $s_2$  which are also inside  $s_1$ 
  for ( $x \in \mathcal{X}_2$ ) {
    | if  $x \in s_1$  then
    | |  $N_{2 \subset 1} \leftarrow N_{2 \subset 1} + 1$ ;
  }
  // Compute approximation of the intersection volume
  Intersection :=  $\frac{V_1 \cdot N_{2 \subset 1} + V_2 \cdot N_{1 \subset 2}}{2N_{\text{samples}}}$ ;
  return  $\frac{\text{Intersection}}{\min(V_1, V_2)}$  ;

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*Sub-optimality of evaluations* Hill-climbing has to evaluate the complete objective function  $\bullet$  each time it considers a primitive, which is particularly costly at the beginning of the algorithm where the solution  $\mathcal{S}_F$  is still empty, since no Cuboid candidate would intersect with it. After selecting a large Cuboid, the set of available, *i.e.* compatible cuboids gets dramatically reduced, hence resulting in an acceleration of the search of Hill-Climbing as shown in Figure 4 of main paper, which however converges to sub-optimal solutions. In contrast, MCTS and our algorithm evaluate the objective function only at the end of an iteration when a complete solution is complete.

### 3.2 Binary-Tree MCTS

In our non-binary tree adaptation of MCTS, the search algorithm spends many iterations on iterating the first levels of the tree which might contain many, mutually incompatible, primitives. Therefore, this can limit the exploitation ca-

**Algorithm 2:** Hill-climbing algorithm

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<b>Result:</b> Set of selected Cuboid $\mathcal{S}_F$ <b>Input:</b> Set of proposal Cuboid $\mathcal{S}$ ; Threshold $\eta$ ; Final solution $\mathcal{S}_F := \emptyset$ ; Current best loss $\ell^* := +\infty$ ; Current best Cuboid $s^* := \emptyset$ ; Set of available Cuboid $\mathcal{S}_A := \mathcal{S}$ ; Evaluations counter $N_{eval} := 0$ ;	<pre> <b>while</b> <math>\mathcal{S}_A \neq \emptyset</math> <b>do</b>   <math>s^* \leftarrow \emptyset</math>;   <b>for</b> ( <math>s \in \mathcal{S}_A</math> ) {     <b>if</b> <math>s.isCompatible(s_f, \eta)</math> <b>then</b>       <math>\mathcal{S}_F.add(s)</math>;       <math>\ell \leftarrow evalObjFunc(\mathcal{S}_F)</math>;       <math>N_{eval} \leftarrow N_{eval} + 1</math>;       <b>if</b> <math>\ell &lt; \ell^*</math> <b>then</b>         <math>\ell^* \leftarrow \ell</math>;         <math>s^* \leftarrow s</math>;       <math>\mathcal{S}_F.remove(s)</math>;     <b>else</b>       <math>\mathcal{S}_A.remove(s)</math>;   }   <math>\mathcal{S}_F.add(s^*)</math>; <b>return</b> <math>\mathcal{S}_F, N_{eval}</math> </pre>
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pabilities of MCTS as the algorithm prioritizes nodes that have not been visited yet. In our binary-tree adaptation of MCTS, every level of the tree corresponds to selecting or skipping a primitive. The resulting tree structure, hence, trades tree-breadth for tree-depth, which enables better exploitation. However, due to a large depth of the tree, MCTS does not explore solutions in the bottom of the tree, hence we observed only minor improvements over its non-binary adaptation. We argue that our MonteBoxFinder method can be seen as an adaptive version of binary-tree MCTS. In contrast to binary-tree MCTS, as we show in Figure 1, the tree equivalent of our method is able to adapt its structure during the search and enable better update mechanism leading to faster convergence.

## 4 More results

### 4.1 Qualitative results

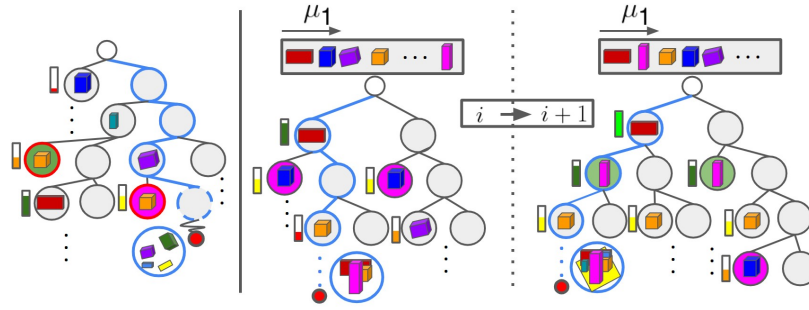
We show more qualitative results in Figure 2.

### 4.2 Progress plots

We show more progress plots in Figure 3.

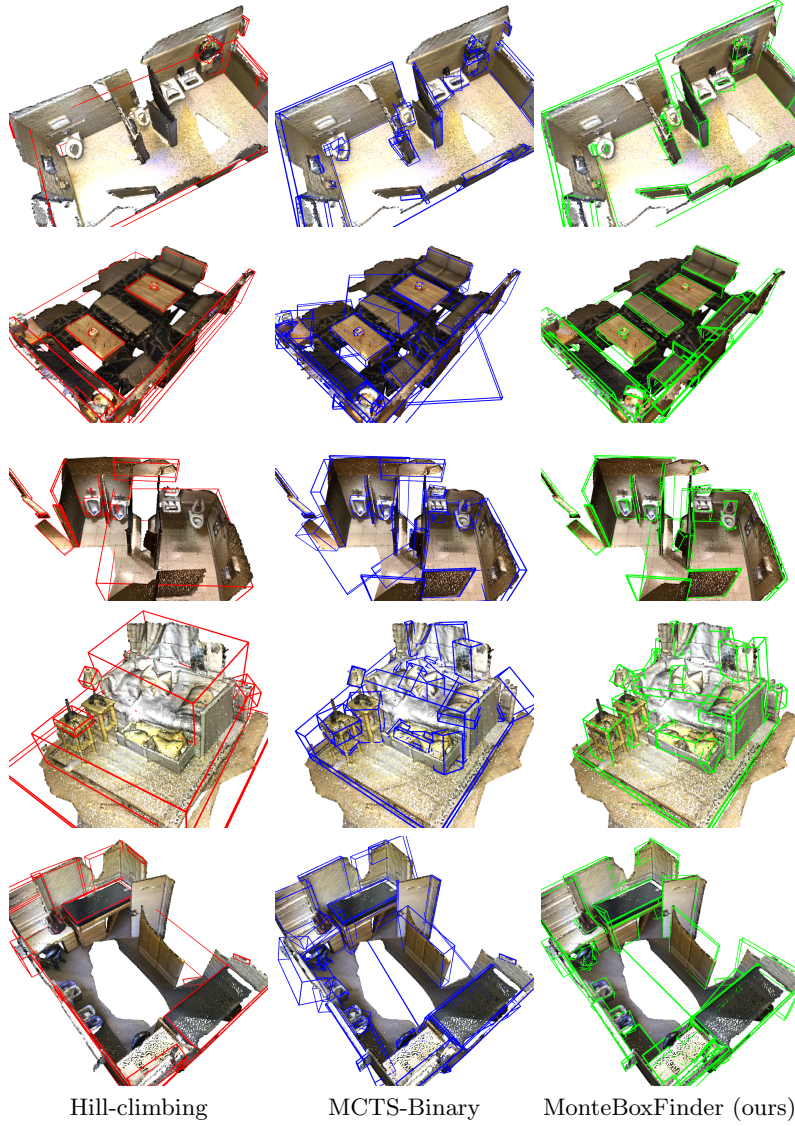
## References

1. Skiena, S.: The Algorithm Design Manual. Springer (2010)

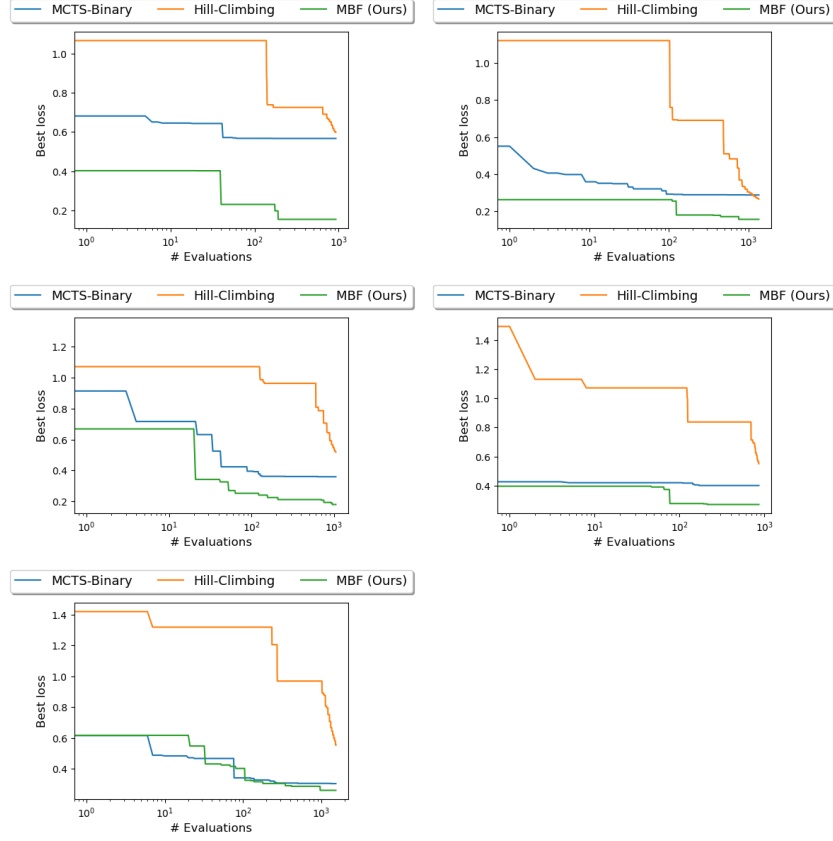


Binary-tree MCTS      2 iterations of our algorithm interpreted with a binary tree

**Fig. 1.** We observe that behavior of our algorithm can be interpreted as an adaptive binary-tree MCTS, even though we do not explicitly define a tree structure. As MCTS is bound by its tree structure, it will invest iterations into exploring primitives in the upper part of the tree, even those with low confidence, visualized as **colored bars**. Further more, as indicated with **colored circles**, MCTS models confidences of same primitives in different parts of the tree independently. **Blue circles** indicate a selection path of a single MCTS iteration that fails to extract meaningful proposals due to aforementioned difficulties. In contrast, our method sorts at each iteration primitives according to their confidences  $\mu_1$  and will focus more easily on more promising primitives. In addition, as indicated by **colored circles** we only model a single confidence per primitive. These features enable our method to converge faster to good solutions in practice.



**Fig. 2. Qualitative results.** Hill-climbing often selects large cuboids that span across multiple different objects. MCTS does better, but sometimes yields outliers (second and fifth row). In contrast, our algorithm outperforms both methods and is able to successfully reconstruct many more details.



**Fig. 3. Samples of progress plots.** Our method consistently outperforms its base-lines, i.e. the Hill-Climbing algorithm and MCTS, as it converges faster to a better solution. These plots correspond to the scene examples in Figure 2